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AN ITERATIVE PROCESS TO SOLVE
THE GRAPH-COLORING PROBLEM

by

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United States Naval Postgraduate School



THESIS

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TO
SOLVE THE GRAPH-COLORING PROBLEM

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October 1969

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An Iterative Process
to
Solve the Graph-Coloring Problem

by

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requirements for the degree of

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ABSTRACT

The intent of this paper is to describe a method of coloring a map and to present an algorithm for the solution of this problem. A computer program was developed to provide solutions to the problem of coloring a map which consists of a finite number of areas. This algorithm may also be applied to problems other than map-coloring.

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TABLE OF SYMBOLS

SYMBOL	MEANING
M	Number of equations used to describe the border constraints for a map.
N	Number of areas to be colored
x_i	Variable representing the color of area i
$A_{M \times N}$	The $M \times N$ matrix defining regions having common borders
\bar{a}_j	The j^{th} column of $A_{M \times N}$
\underline{a}_i	The i^{th} row of $A_{M \times N}$
$A'_{M \times N}$	$A_{M \times N} + [\bar{a}_j, \dots, \bar{a}_j]_{1 \times N}$
β_j	The number of zero elements of each column \bar{a}_j that are not associated with the constraint equations
α_j	The number of zero elements of each column \bar{a}_j

I. INTRODUCTION

When coloring a geographical map, it is customary to use different colors for adjacent regions - where adjacent means a common finite boundary. It has been conjectured, but not proved, that "For any subdivision of the plane into non-overlapping regions, it is always possible to mark the regions with one of the numbers 0, 1, 2, 3, in such a way that no two adjacent regions receive the same number."¹

The statement that a map can be colored with four colors is accredited to Moebius, who first proposed it in 1840. This conjecture has never been proven; however, a satisfactory counter-example has not been demonstrated either. A proof for five-coloring a map may be found in [Courant and Robbins 1941] and [Oystein 1967].

This thesis will use the four-coloring conjecture without proof and present an algorithm that utilizes an iterative routine to determine what color each region should be. A similar iterative technique for general integer programming problems is demonstrated in [Greenberg April 1969] and [Greenberg May 1969], which present the proofs. We are interested in showing whether the iterative technique can be used for large-scale problems such as the map-coloring problem.

The problem formulation suggested by Gomory [Dantzig 1960] is used. Let the map regions be denoted by $i = 1, 2, \dots, N$ and let x_i be an integer valued variable such that

$$0 \leq x_i \leq 3 \quad (1)$$

¹Courant, Richard and Robbins, Herbert; What is Mathematics?, p. 247, Oxford University Press, 1941.

where the four values $x_i = 0, 1, 2, \text{ or } 3$ correspond to four different colors. Since any two adjacent areas must have different colors, the constraint may be written in an either/or form

$$\text{either } x_i - x_j \geq 1 \quad \text{or} \quad x_j - x_i \geq 1 \quad (2)$$

if i and j have a common border. Equation (2) may be rewritten as:

$$\begin{aligned} x_i - x_j &\geq 1 - 4\delta_{ij} & (\delta_{ij} = 0, 1) \\ x_j - x_i &\geq -3 + 4\delta_{ij} \end{aligned} \quad (3)$$

Instead of using (2) and (3) we can simplify the problem by using

$$x_i - x_j \neq 0 \quad (4)$$

where i and j have a common border.

This formulation can be written in matrix form

$$Ax \neq 0 \quad (5)$$

where the elements of $A_{M \times N}$, a_{ij} , are either 0, -1, or 1. We can now enumerate all values of the left side of (5). The column with the least number of zeroes is now selected from $A_{M \times N}$, \bar{a}_j ,² and the value of the corresponding x_j is set equal to one. The column \bar{a}_j is then added column by column to the matrix $A_{M \times N}$ to form a new matrix. The process continues in the same way by searching the new matrix for the column with the least number of zeros. In each iteration the corresponding x_j is increased by one. The column values represent the value of the left side of (5). A solution occurs when a column is achieved that has no zeros.

²The speed with which the algorithm reaches a solution is increased by modifying the method of choosing \bar{a}_j . This modification is explained in Section II. A. 4.

The algorithm presented is based upon the four-color conjecture, but is not restricted to the use of only four colors, if more colors are desired.

II. THE COLORING PROCESS

A. THE ALGORITHM

1. Let $x_i = 0$ for $i = 1, \dots, N$.
2. For every 2 regions of the map with a common border it is required that $x_i - x_j \neq 0$, $i, j, = 1, \dots, N$; $i \neq j$.
3. Form the $M \times N$ matrix ($A_{M \times N}$) defined in step 2, where M is the number of equations required to describe the border constraints and N is the number of areas to be colored. The elements of $A_{M \times N}$, a_{ij} , are either -1, 0, or 1. Note:
 - a. For each row, \underline{a}_i , $i = 1, \dots, M$, there are only two non zero elements, -1, and 1, corresponding to the constraints in step 2.
 - b. For column \bar{a}_j , $j = 1, \dots, N$, has at least one non zero element and fewer than M non zero elements.
4. Let α_j , $j = 1, \dots, N$, be the number of zero elements of each \bar{a}_j . Let β_j , $j = 1, \dots, N$, be the number of zero elements of each \bar{a}_j which appear in those positions where $a_{ij} = 0$. Note: $\beta_j \leq \alpha_j$, $j = 1, \dots, N$. Find the column of $A_{M \times N}$ with the smallest β_j . Let j^{**} be the smallest j such that $\beta_{j^{**}} = \min_j \{\beta_j\}$. Let $j^* = j^{**}$.
5. Let $x'_{j^*} = x_{j^*} + 1$. Check all the constraints in which x'_{j^*} appears to insure that the constraint equations (4) are satisfied. If the constraint equations (4) are not satisfied, continue increasing x_{j^*} in increments of one until either (4) is satisfied or $x'_{j^*} = 3$. If $x'_{j^*} = 3$ and (4) is not satisfied, set $x'_{j^*} = 0$ and let j^* be the

smallest $j > j^{**}$ such that $\beta_{j^*} = \min_j \{\beta_j\}$ and repeat step 5. If no such j^* exists, then find the smallest α_j and proceed as with the β_j 's. If no α_j exists such that x_{j^*} and x_i satisfy (4) for all $i \neq j^*$, go to step 9.

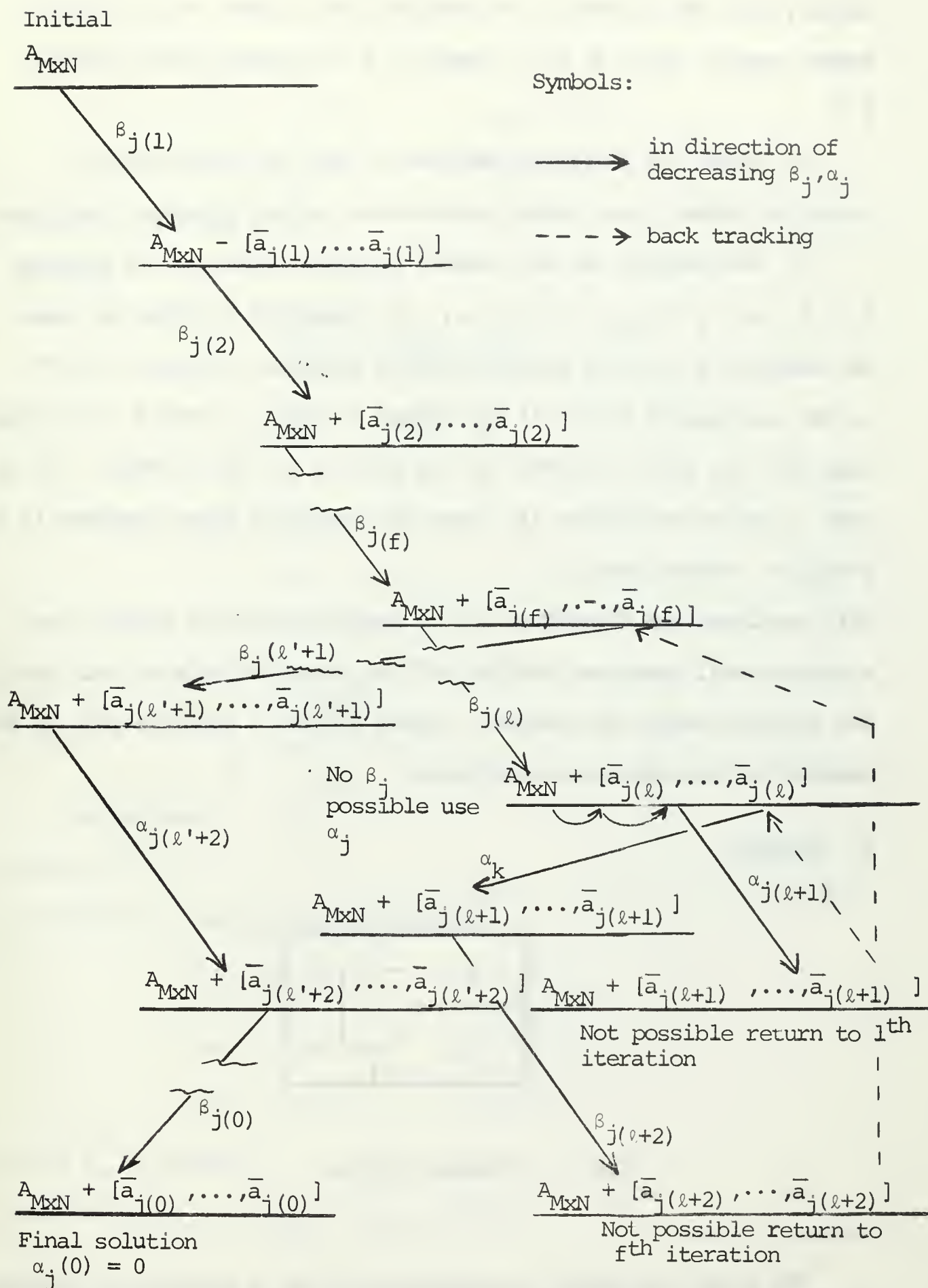
6. Determine the elements of column \bar{a}_{j^*} by substituting the presently defined values of x_i , $i = 1, \dots, N$, into the constraint equations (4). Add the column \bar{a}_{j^*} , $j = 1, \dots, N$, to each column of $A_{M \times N}$. Let $A'_{M \times N} = A_{M \times N} + [\bar{a}_{j^*}, \dots, \bar{a}_{j^*}]$.

7. Let α_j , $j = 1, \dots, N$, be the number of zero elements of each \bar{a}'_j . Let β_j , $j = 1, \dots, N$, be the number of zero elements of each \bar{a}'_j which appear in those positions where $a_{ij} = 0$. Let j^{**} be the smallest j such that $\beta_j = \min_j \{\beta_j\}$. Let $j^* = j^{**}$.

8. Continue steps 5 through 8 until $\alpha_j = 0$, in which case a solution has been reached. Note: If $x_j = 3$ at any point during steps 5 through 8, disregard the corresponding column \bar{a}_j in any future iteration, since the corresponding x_j is at its maximum possible value.

9. Let the column \bar{a}_{j^*} be the one used to generate the latest $A'_{M \times N}$ matrix. For each zero element of \bar{a}_{j^*} there are 2 regions, g and h , with a common border where $x_g - x_h = 0$, i.e., a violation of (4). Let F be the set of all such regions g and h . Denote by the subscript f the element of F .

10. For each f , increment the x_f 's such that $x_f = x_f + 1$. Check all of the constraints in which x_f appears to see if the constraint equations (4) are satisfied. If for any f , an element of F , the constraint equations (4) are satisfied set $f = j^*$ and go to step 6. If (4) is not satisfied for any i , not an element of F , where regions i and f have a common



Possible Path to a solution to the map-coloring problem

Figure 1

border, note the iteration that defined the present value of x_i .

Repeat step 10 until $x_f = 3$. Reset $x_f = 0$. Repeat step 10 for all F .³

11. Order the iterations defined in step 10 in decreasing iteration number. Let these iterations be called possible iterations.

12. Reconstruct the most recent possible iteration and examine $\beta_j = \beta_{j^*}$ or $\alpha_j = \alpha_{j^*}$, $j = j^* + 1, \dots, N$, depending on which was used to determine x_{j^*} in the iteration being examined. Choose a new j^* in the same manner as the j^* was chosen in step 5. When a j^* is found such that x_{j^*} and x_i satisfy (4) for all $i \neq j^*$, go to step 6. If no such j^* exists satisfying (4), then the iteration being examined is not possible. Repeat step 12.

This completes the algorithm. If no possible solution exists, the algorithm will continue "backing up" the generated matrix chain until the original matrix is reached. Figure 1 shows a possible path to the solution of the map-coloring problem.

B. EXAMPLE

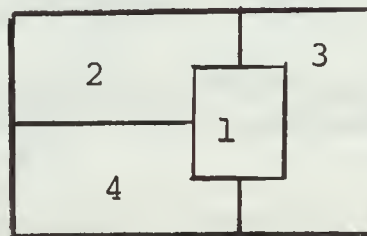


Fig. 2. Example Problem

³The speed with which the algorithm reaches a solution is increased by starting the backtracking procedure when no variable associated with the minimum β_j or α_j can be set at any integer value between 1 and 3.

where

$$\begin{array}{ll} x_1 - x_2 \neq 0 & x_2 - x_3 \neq 0 \\ x_1 - x_3 \neq 0 & x_2 - x_4 \neq 0 \\ x_1 - x_4 \neq 0 & x_3 - x_4 \neq 0 \end{array}$$

set $x_i = 0$ for all i

we write

	x_1	x_2	x_3	x_4
	1	-1	0	0
	1	0	-1	0
	1	0	0	-1
	0	1	-1	0
	0	1	0	-1
	0	0	1	-1
α_j	3	3	3	3
β_j	3	3	3	3

column used *

then $x_1 = 1$

the column to add (\bar{a}_{j*}) is defined by:

$$\begin{array}{ll} x_1 - x_2 = 1 & x_2 - x_3 = 0 \\ x_1 - x_3 = 1 & x_2 - x_4 = 0 \\ x_1 - x_4 = 1 & x_3 - x_4 = 0 \end{array}$$

where $\bar{a}_{j*}^T = [111000]$

then adding $[\bar{a}_{j*}, \dots, \bar{a}_{j*}] + A_{M \times N}$ we write

	x_1	x_2	x_3	x_4
	2	0	1	1
	2	1	0	1
	2	1	1	0
	0	1	-1	0
	0	1	0	-1
	0	0	1	-1
α_j	3	2	2	2
β_j	3	1	1	1

column used *

then $x_2 = 2$

the column to add (\bar{a}_{j*}) is defined by:

$$\begin{array}{ll} x_1 - x_2 = -1 & x_2 - x_3 = 2 \\ x_1 - x_3 = 1 & x_2 - x_4 = 2 \\ x_1 - x_4 = 1 & x_3 - x_4 = 0 \end{array}$$

where $\bar{a}_{j*}^T = [-111220]$

then adding $[\bar{a}_{j*}, \dots, \bar{a}_{j*}] + A_{M \times N}$ we write

	x_1	x_1	x_3	x_4
	0	-2	-1	-1
	2	1	0	1
	2	1	1	0
	2	3	1	2
	2	3	2	1
	0	0	1	-1
α_j	2	1	1	1
β_j	1	1	0	0

column used *

then $x_3 = 3$

the column to add (\bar{a}_{j*}) is defined by:

$$\begin{array}{ll} x_1 - x_2 = -1 & x_2 - x_3 = -1 \\ x_1 - x_3 = -2 & x_2 - x_4 = 2 \\ x_1 - x_4 = 1 & x_3 - x_4 = 3 \end{array}$$

where $\bar{a}_{j*}^T = [-1 \ -2 \ 1 \ -1 \ 2 \ 3]$

then adding $[\bar{a}_{j*}, \dots, \bar{a}_{j*}] + A_{M \times N}$ we write

$$\begin{array}{cccc} & x_1 & x_2 & x_3 & x_4 \\ & \hline & 0 & -2 & -1 & \\ -1 & & -2 & -3 & \\ 2 & & 1 & 1 & \\ -1 & & 0 & -2 & \\ 2 & & 3 & 2 & \\ 3 & & 3 & 4 & \\ \hline \alpha_j & 1 & 1 & 0 & \\ & & & * & \end{array}$$

there is no need to calculate column \bar{a}'_4 because there are no zeros in column \bar{a}'_3 and a solution has been reached. The solution is:

Area number	Color number
1	1
2	2
3	3
4	0

III. CONCLUSION

The algorithm described will solve the map-coloring problem. Although only four colors were used in the example in Section II-B, the algorithm can be adapted to solve a problem for any feasible number of colors. Several representative map-coloring problems were solved using the enclosed computer program. Appendix A shows a solution to four-coloring the continental United States. In the representative problems solved, both four and five colors were used.

If the solution is not feasible, e.g., the use of four colors when five colors are required, as in the split-state problem, the algorithm may consume several hours of computer time before showing that no solution exists. The computer time is consumed in back-tracking the solution path and in investigating every possible solution path for both the α 's and the β 's.

APPENDIX A

SAMPLE PROBLEM SOLUTION

The solution for four-coloring the continental United States of America, including boundaries and major water bodies, is as follows:

STATE	COLOR
Maine	0
New Hampshire	3
Vermont	0
Massachusetts	2
Rhode Island	3
Connecticut	0
New York	3
New Jersey	0
Pennsylvania	2
Delaware	3
Maryland	0
Virginia	2
North Carolina	3
South Carolina	0
Georgia	2
Florida	0
Alabama	3
Tennessee	1
West Virginia	1
Ohio	0
Indiana	2
Kentucky	3
Illinois	0
Wisconsin	2
Michigan	3
Minnesota	0
Iowa	3
Missouri	2
Arkansas	3
Mississippi	2
Louisiana	0
Texas	2
Oklahoma	0
Kansas	3
Nebraska	0
South Dakota	1
North Dakota	3
Montana	0
Wyoming	2

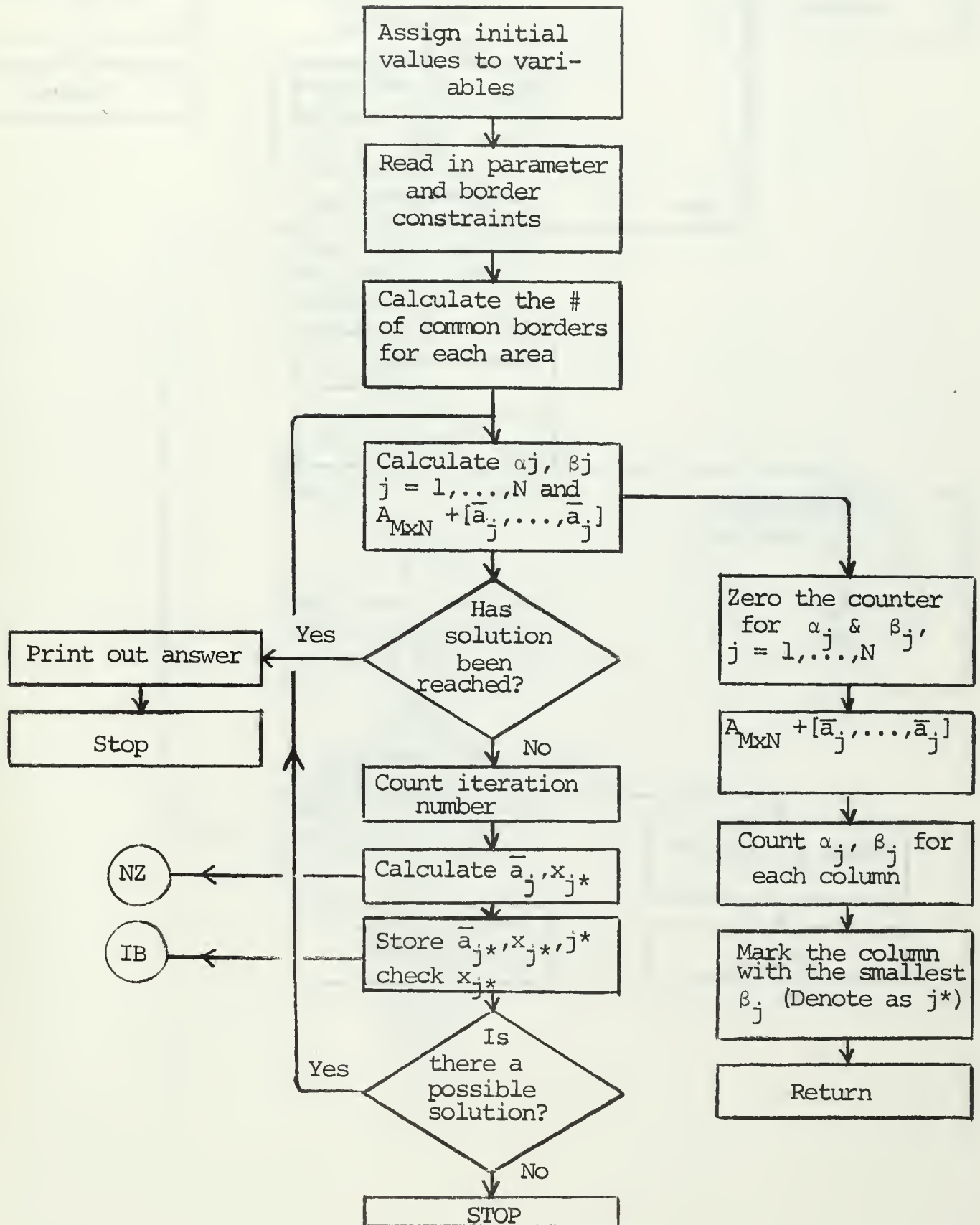
Colorado	1
New Mexico	3
Arizona	1
Utah	3
Idaho	1
Washington	0
Oregon	2
Nevada	0
California	3

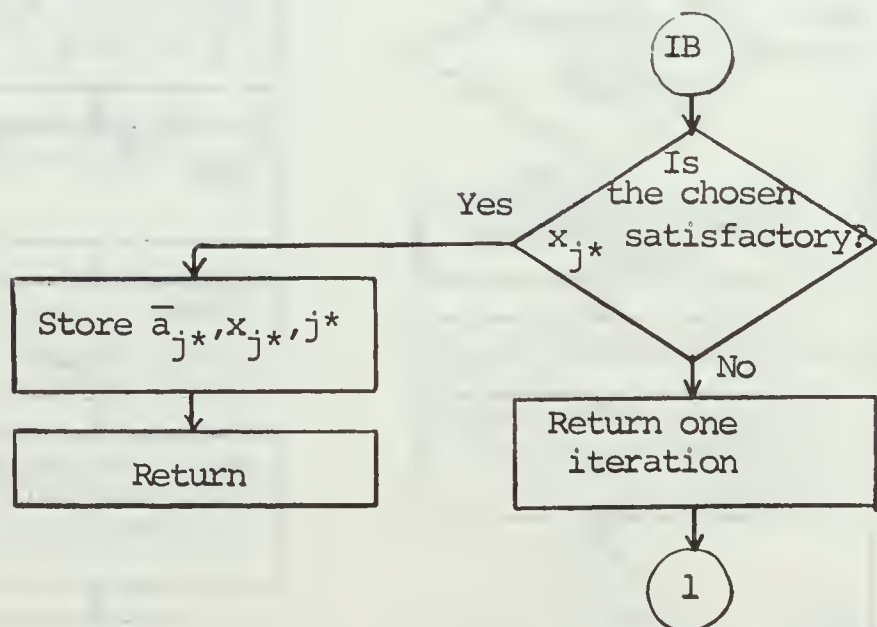
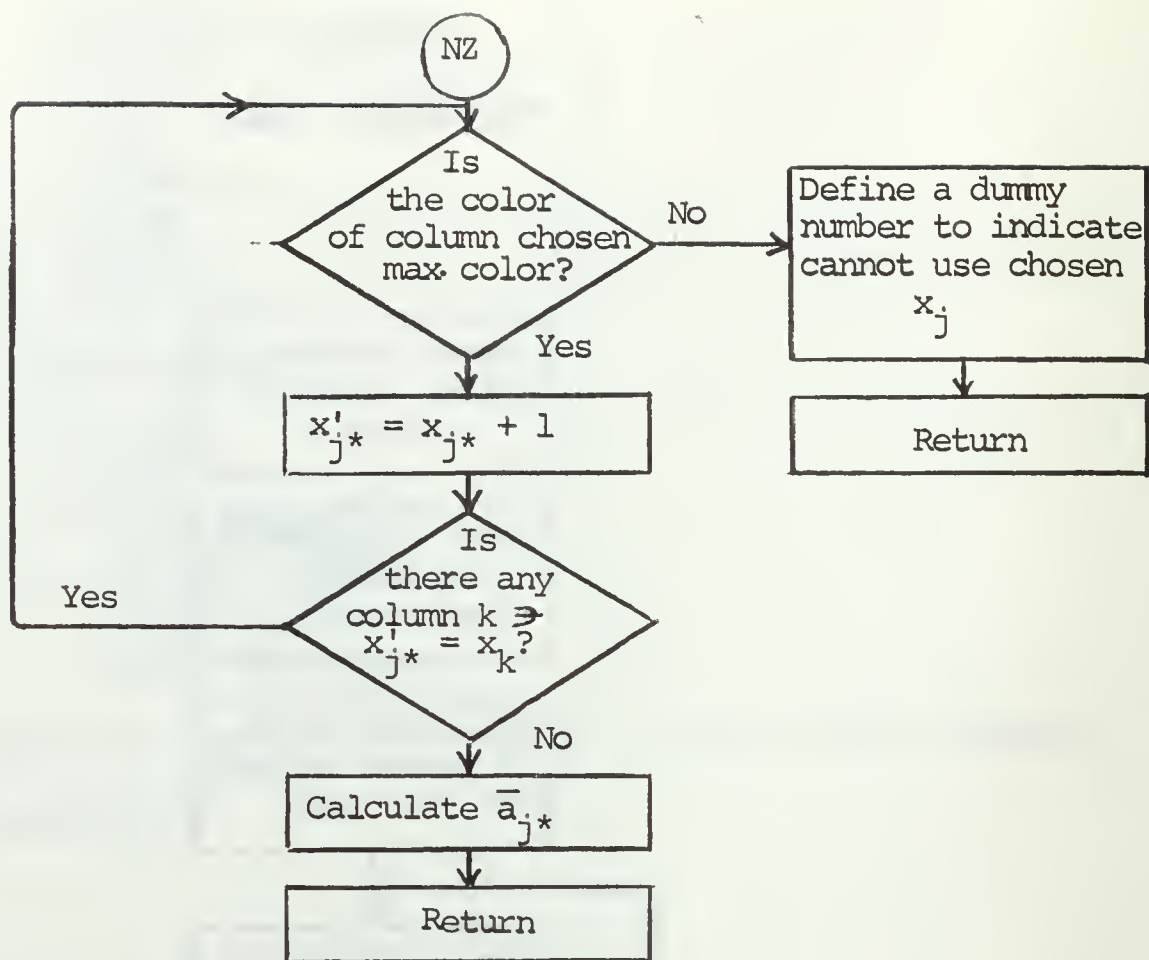
BORDER AREAS

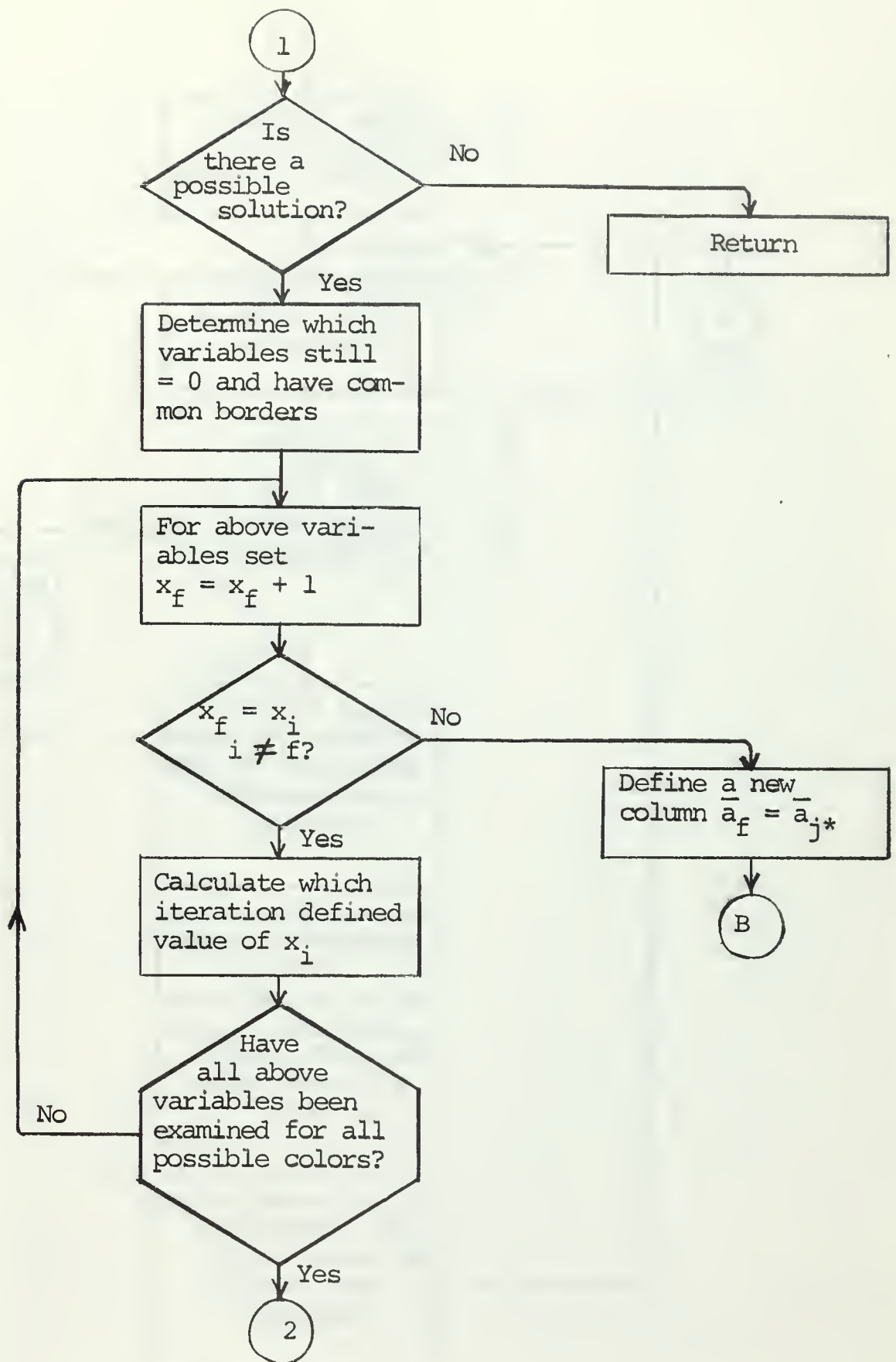
Canada	2
Mexico	0
Water - including:	1
Atlantic Ocean	
Pacific Ocean	
Great Lakes	
Great Salt Lake	

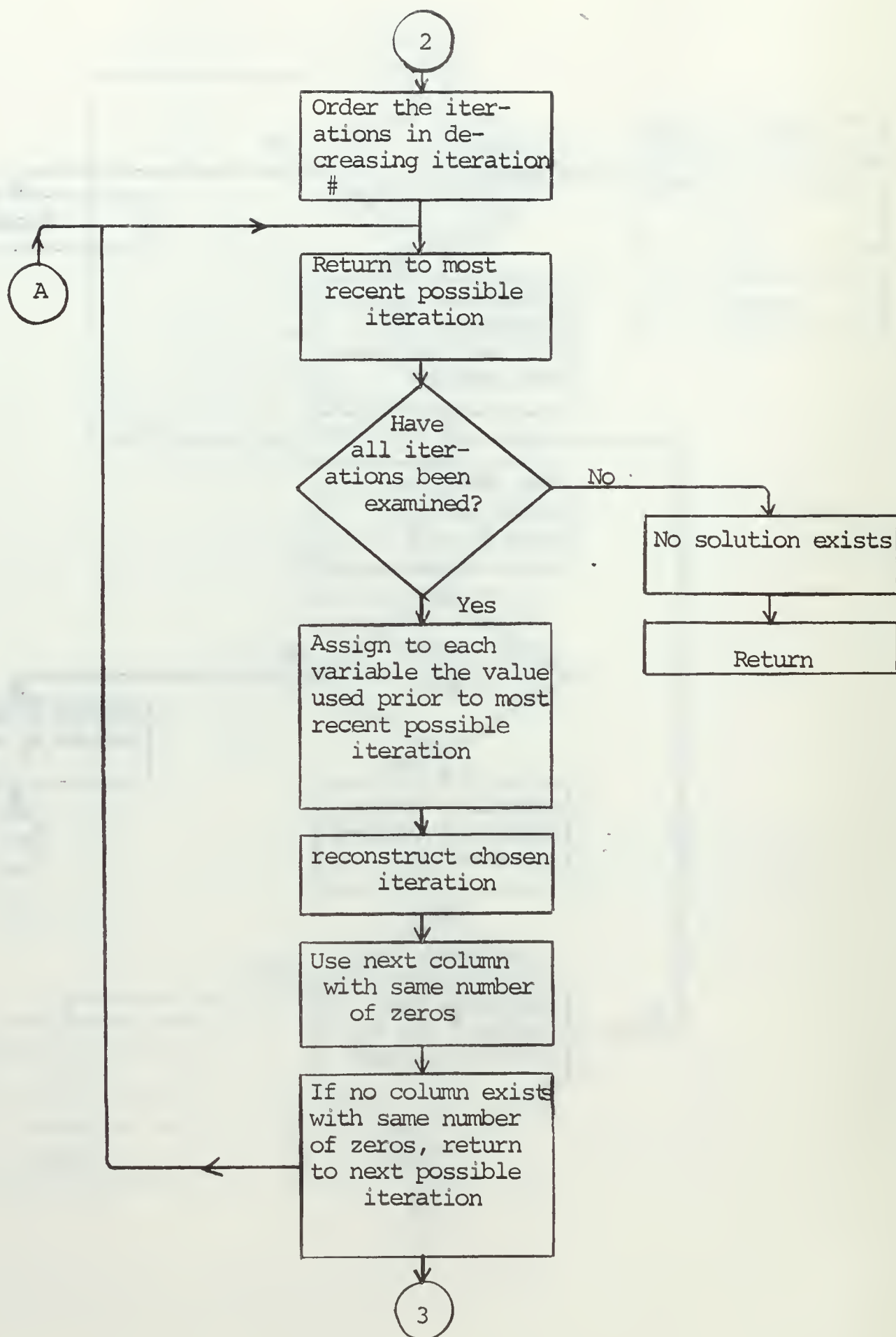
The above problem was solved in 97.2 seconds on the IBM 360/67 computer and used 81 iterations.

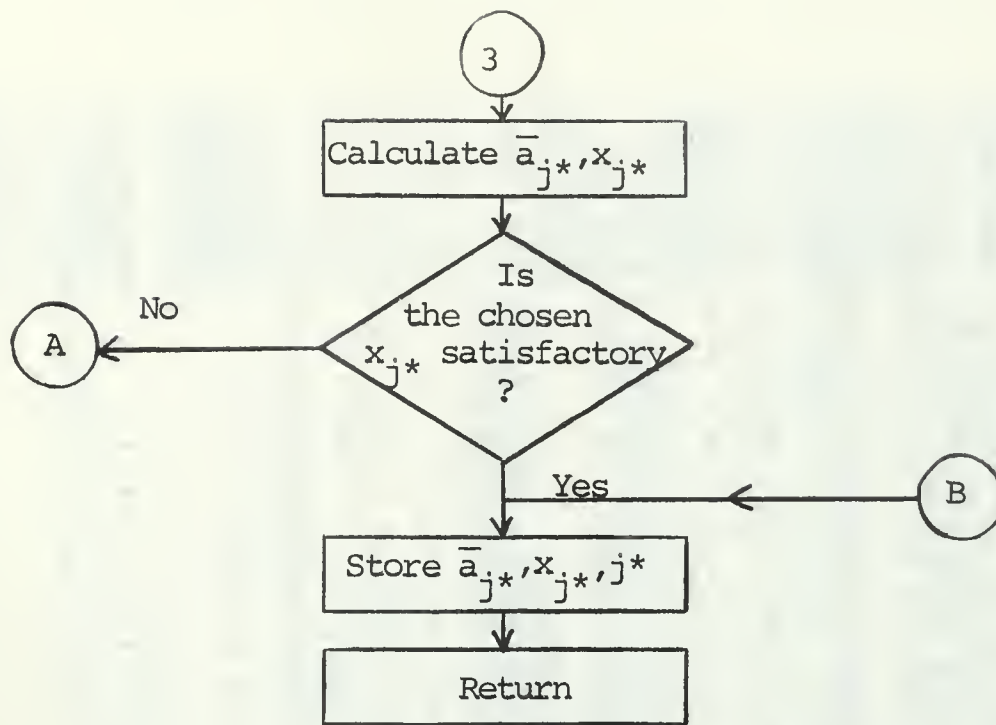
APPENDIX B
PROGRAM FLOW CHART












```

110 MV(I)=0
115 DO 115 I=1,N
115 IADD(I)=0
117 DO 117 I=1,N
117 DO 117 J=1,NN
117 ITAB(I,J)=0
C READ IN THE INITIAL VALUES DEFINING BORDER AREAS
119 READ(5,120) I,J,K,ITAB(I,J),ITAB(I,K),IDUM
120 FORMAT(6I5)
C IF(IDUM.EQ.0) GO TO 119
C DETERMINE THE NUMBER OF BORDERING AREAS FOR EACH STATE
DO 121 J=1,NN
DO 121 I=1,N
IJ=ITAB(I,J)
121 KOLM(J)=KOLM(J)+IABS(IJ)
C DETERMINE THE COLUMN TO USE FOR THE NEXT ITERATION AND MARK IT
125 CALL NMTX(NN,N,ITAB,IADD,MAT,II,DUM,NZ,ITSV,NZCON)
C IF SOLUTION HAS BEEN REACHED, PRINT OUT THE ANSWER
C IF(DUM.GT.0.0) GO TO 130
C INCREASE COUNTER BY ONE
KAT=KAT+1
C DETERMINE THE NEW VALUE OF THE VARIABLE, INDICATE IF SOLUTION IS
C INFEASIBLE SO THAT A NEW PATH CAN BE TAKEN, AND DEFINE THE NEW COLUMN
C TO BE ADDED TO THE INITIAL TABLEAU
140 CALL NZERO(NN,II,ITAB,ITSV,MV,E,KLR,IADD,BA,KOLM)
C IF(BA.NE.0.0) GO TO 140
C SAVE THE COLUMN ADDED TO THE INITIAL TABLEAU/VARIABLE USED. IF SOLUTION
C FEASIBLE, GO BACK TO FIRST FEASIBLE TABLEAU AND PICK A NEW VARIABLE TO
CALL IBACK(NN,IADD,KNT,II,NZ,ITAB,ITSV,DUM,MV,MAT,E,KLR,KOLM,
INZCON,NP3,NNX3)
IF(DUM.GT.0.0) GO TO 130
GO TO 125
C WRITE OUT COLORS FOR EACH AREA.
130 DO 135 I=1,NN
WRITE(6,150) I,MV(I)
150 FORMAT(10X,'FOR AREA NUMBER USE COLOR NUMBER',T27,I3,T48,I3)
135 CONTINUE
STOP
END

SUBROUTINE NMTX(NN,N,IM,IT,MT,II,D,NZ,ISV,NZCON)
SUBPROGRAM TO DO THE FOLLOWING:
(1) ADD COLUMN WITH LEAST NUMBER OF ZEROS TO INITIAL MATRIX
(2) IF ANY COLUMN HAS NO ZEROS STOP THE ITERATION
SYMBOLS ARE:
MT=MAT; IM=ITAB,D=DUM, IT=IADD, ISV=ITSV

```

MAIN0058
 MAIN0060
 MAIN0080
 MAIN0090
 MAIN0095
 MAIN0096
 MAIN0097
 MAIN0100
 MAIN0110
 MAIN0111
 MAIN0115
 MAIN0116
 MAIN0117
 MAIN0118
 MAIN0120
 MAIN0130
 MAIN0142
 MAIN0150
 MAIN0153
 MAIN0155
 MAIN0156
 MAIN0158
 MAIN0160
 MAIN0170
 MAIN0180
 MAIN0190
 MAIN0200
 MAIN0210

NMAT0010

```

C      II: THE SUBSCRIPT OF THE COLUMN WITH LEAST NUMBER OF ZEROS
C      NZ: THE NUMBER OF ZEROS IN EACH COLUMN
C      NZCON: COUNTS THE NUMBER OF ZEROS IN A COLUMN THAT ARE NOT DEFINED IN THE
C      BORDER EQUATIONS
C
C      DIMENSION MT(N,NN),IM(N,NN),IT(N),NZ(NN),ISV(NN),NZCON(NN)
C      ZERO THE VALUE OF THE NUMBER OF ZEROS IN EACH COLUMN
C      DO 320 I=1,NN
C      NZCON(I)=0
C      320 NZ(I)=0
C      ADD THE DESIGNATED COLUMN TO THE INITIAL MATRIX AND COUNT THE NUMBER
C      OF ZEROS IN EACH COLUMN
C      DO 300 J=1,NN
C      DO 305 K=1,N
C      MT(K,J)=IM(K,J)+IT(K)
C      IF(MT(K,J).EQ.0) CALL KONFLT(NN,J,K,NZ,NZCON,IM,N)
C      305 CONTINUE
C      IF ANY COLUMN HAS NO ZEROS, A SOLUTION HAS BEEN REACHED
C      IF(NZ(J).EQ.0) GO TO 350
C      IF VARIABLE IS AT MAXIMUM VALUE, DEFINE NUMBER OF ZEROS TO BE AT MAXIMUM SO
C      THAT COLUMN WILL NOT BE FURTHER CONSIDERED
C      IF(ISV(J).EQ.N) NZ(J)=N
C      300 CONTINUE
C      FIND COLUMN WITH LEAST NUMBER OF ZEROS
C      CALL IVAR(NN,NZ,II,NZCCN,N)
C      RETURN
C      A SOLUTION HAS BEEN REACHED
C      350 D=NN
C      II=J
C      WRITE(6,390) II
C      390 FORMAT(//10X,'LAST VARIABLE USED IS',I35,I3)
C      RETURN
C      END
C
C      SUBROUTINE IVAR(NN,MZ,II,NC,N)
C      SUBROUTINE TO DETERMINE WHICH COLUMN HAS THE LEAST NUMBER OF ZEROS THAT
C      ARE NOT ASSOCIATED WITH BORDER AREAS.
C      SYMBOLS ARE:
C      MZ=NZ; NC=NZCON
C      NP: DUMMY TO INSURE ALL COLUMNS ARE CHECKED
C      II: DENOTES THE COLUMN(VARIABLE) TO BE USED.
C
C      DIMENSION MZ(NN),NC(NN)
C      ZERO VALUES
C      II=1
C      NP=NN-1
C      IF AN AREA IS AT ITS MAXIMUM COLOR, INSURE IT IS NOT CHOSEN

```

NMAT0020
 NMAT0023
 NMAT0024
 NMAT0025
 OF
 NMAT0030
 NMAT0040
 NMAT0050
 NMAT0060
 NMAT0065
 NMAT0080
 MAXIMUM SO
 NMAT0085
 NMAT0090
 NMAT0100
 NMAT0130
 NMAT0140
 NMAT0142
 NMAT0144
 NMAT0145
 NMAT0150
 NMAT0160

IVAR0020
 IVAR0030
 IVAR0035


```

C      DO 210 I=1,NN
C      210 IF(MZ(I).EQ.N) NC(I)=N
C      PICK THE COLUMN WITH THE LEAST NUMBER OF NONBORDERING ZERO AREAS
C      DO 200 K=1,NP
C      200 IF(NC(K+1).LT.NC(II)) II=K+1
C      RETURN
C      END

```

IVAR00037
IVAR00038

IVAR00040
IVAR00050
IVAR00060
IVAR00070

```

C      SUBROUTINE NZERO(N,NN,II,IT,ITSV,MV,E,KL,IA,BA,KM)
C      A SUBPROGRAM TO DETERMINE WHICH EQUATIONS CORRESPOND TO THE VARIABLE
C      CHOSEN. // SYMBOLS ARE AS DEFINED IN MAIN PROGRAM; OTHER SYMBOLS ARE:
C      ITSV: DENOTES VARIABLE IS AT MAXIMUM VALUE
C      BA: DUMMY TO INDICATE THAT 2 ADJACENT AREAS HAVE THE SAME COLOR
C      K2: COUNTER TO INDICATE WHEN ALL BORDER AREAS HAVE BEEN EXAMINED
C      E: A DUMMY TO INDICATE THAT TWO VARIABLES ARE EQUAL AND AT MAX
C      OTHER SYMBOLS: IT=ITAB; KL=KLR
C
C      SUBROUTINE NZERO(N,NN,II,IT,ITSV,MV,E,KL,IA,BA,KM)
C      NZERO010

```

```

C      DIMENSION IT(N,NN),ITSV(NN),MV(NN),IA(N),KM(NN)
C      ZERO THE COUNTERS AND DUMMIES
C      BA=0.0
C      E=0.0
C      K2=0
C      IF COLUMN CHOSEN IS AT MAXIMUM VALUE, RETURN AND CHOOSE A NEW COLUMN
C      IF(MV(II).EQ.KL) GO TO 610
C      INCREASE VALUE OF VARIABLE BY ONE
C      MV(II)=MV(II)+1
C      IF COLOR IS AT THE MAXIMUM VALUE, MAKE A NOTE OF THE VARIABLE
C      IF(MV(II).EQ.KL) ITSV(II)=N
C      CHECK EACH ROW FOR VARIABLE USED TO CHECK THE COLOR OF THE ADJOINING
C      DO 600 JJ=1,N
C      IF(IT(JJ,II).NE.0) CALL INFEAS(N,NN,MV,IT,II,JJ,KL,IA,BA,KM,K2)
C      IF(BA.NE.0.0) RETURN
C      600 CONTINUE
C      RETURN
C      610 E=N
C      END

```

NZERO0020
NZERO0024
NZERO0025
NZERO0026

NZERO0035

NZERO0036

NZERO0037
AREAS
NZERO0038
NZERO0040
NZERO0045
NZERO0050
NZERO0060
NZERO0064
NZERO0065
NZERO0070

```

C      SUBROUTINE IBACK(N,NN,IA,KNT,II,NZ,IN,ITSV,D,MV,MT,E,KL,KM,NC,N3,
C      INX3)
C      SUBPROGRAM TO DO THE FOLLOWING:
C      (1) SAVE THE COLUMN ADDED TO THE INITIAL MATRIX
C      (2) SAVE THE INDEX VALUE OF THE COLUMN ADDED TO THE INITIAL MATRIX AND SAVE
C      THE INDICATOR FROM WHERE THE COLUMN CAME (ALPHA OR BETA) AND SAVE THE
C      COLOR OF THE VARIABLE
C      (3) REINDEX THE COLUMN, IF NO OTHER COLUMN CAN BE CHOSEN

```

IBAK00010
IBAK00011


```

C RETURN ONE ITERATION
C 410 KNT=KNT-1
C IF ALL ITERATIONS HAVE BEEN EXAMINED, NO SOLUTION EXISTS
C IF (KNT.EQ.0) GO TO 450
C IF THE VARIABLE WAS AT MAXIMUM VALUE, REZERO THE INDICATOR
C IF (MV(I1).EQ.KL) ITSU(I1)=0
C ZERO THE COUNTERS AND INDICATORS
C L4=0
C L2=0
C KN1=0
C DETERMINE THE COLUMNS ASSOCIATED WITH THE ZERO VALUES IN IADD
C DO 482 I=1,N
C 482 IF (IA(I).EQ.0) CALL KNTVAR(N,NN,IN,I,IITMP,L2)
C IF THE COLOR OF THE ADJACENT AREAS THAT ARE PRESENTLY AT A ZERO COLOR LEVEL
C WERE INCREASED TO THE MAXIMUM COLOR, DETERMINE WHAT AREAS WOULD THEN HAVE
C MATCHING COLORS
C DO 460 J=1,L2
C ASSIGN A TEMPORARY VALUE TO VARIABLE THAT IS AT ZERO LEVEL
C IIS=IITMP(J)
C ASSIGN AN INCREASING COLOR VALUE TO ZERO LEVEL VARIABLES
C DO 461 I=1,KL
C MV(IIS)=I
C ASSIGN INITIAL VALUE TO AN INDICATOR TO SHOW WHEN A COLOR ASSIGNED TO AN
C AREA IS NOT THE SAME AS SOME BORDERING AREA
C K2=KN1
C CHECK EACH ROW FOR THE COLOR USED TO CHECK THE COLOR OF THE ADJOINING
C AREAS.
C DO 462 JJ=1,N
C 462 IF (IN(JJ,IIS).NE.0) CALL LKINF(N,NN,IIS,KN1,IN,MV,MVMCH,JJ)
C IF PRESENT COLOR DOES NOT MATCH WITH ANY BORDER AREAS, STOP THE EXAMINING
C PROCESS AND RETURN TO MAIN PROGRAM.
C IF (KN2.EQ.KN1) CALL NOMCH(N,NN,IIS,II,I,MV,IA,IN,NU,KNT,KL,ITSV,
C 1NX3)
C IF (KN2.EQ.KN1) GO TO 441
C 461 CONTINUE
C REZERO VARIABLE EXAMINED
C MV(IIS)=0
C 460 CONTINUE
C DETERMINE FROM WHICH ITERATION EACH OF THE VARIABLES PICKED ABOVE CAME
C DO 471 J=1,KN1
C ITER(J)=0
C DO 471 I=1,KNT
C IF (MVMCH(J).EQ.KOLSV(N+1,I)) ITER(J)=I
C 471 CONTINUE
C ORDER THE ITERATIONS
C CALL IORDER(N,KN1,ITER)

```

IBAK0100

IBAK0110

IBAK0090

IBAK0091

IBAK0092

IBAK0093

IBAK0094

IBAK0095

IBAK0096

IBAK0097

IBAK0098

IBAK0099

IBAK0100

IBAK0101

IBAK0102

IBAK0103

IBAK0103

IBAK0104

IBAK0105

IBAK0106

IBAK0107

IBAK0108

IBAK0109

IBAK0110

IBAK0111

IBAK0112

IBAK0113

```

C RETURN TO THE MOST RECENT ITERATION DEFINED BY IORDER
411 L4=L4+1
C KNT=ITER(L4)
C IF ALL ITERATIONS HAVE BEEN EXAMINED NO SOLUTION EXISTS
  IF(KNT.LE.C) GO TO 450
C AD=0.0
C ASSIGN TO THE VARIABLES THE VALUES THAT WERE ASSIGNED IN ITERATIONS
  C TO THE CHOSEN ITERATION
  LNT=KNT-1
  DO 475 J=1,NN
    ITSV(J)=0
    MV(J)=0
    DO 476 I=1,LNT
      IF(KOLSV(N+1,I).EQ.J) MV(J)=KOLSV(N+3,I)
      IF(MV(J).EQ.KL) ITSV(J)=N
    C RETRIEVE THE VARIABLE USED IN CHOSEN ITERATION.
    II=KOLSV(N+1,KNT)
    IF(KOLSV(N+1,KNT).EQ.NN) GO TO 411
    C RECONSTRUCT CHOSEN ITERATION MATRIX(MAT)
    DO 425 I=1,NN
      NC(I)=0
      NZ(I)=0
    DO 420 J=1,NN
      DO 419 K=1,N
        MT(K,J)=IN(K,J)+KOLSV(K,LNT)
        IF(MT(K,J).EQ.0) CALL KONFLT(NN,J,K,NZ,NC,IN,N)
    419 CONTINUE
    IF(ITSV(J).EQ.N) NZ(J)=N
    420 CONTINUE
  C USE NEXT COLUMN WITH SAME NUMBER OF ZEROS, IF NO SUCH COLUMN EXISTS RETURN
  C TO NEXT FEASIBLE ITERATION
  IND=KOLSV(N+2,KNT)
  CALL KVAR(NN,NZ,II,AD,MV,ITSV,NC,IND,NU,KNT,NX3)
  KOLSV(N+2,KNT)=IND
  IF(AD.NE.0.0) GO TO 411

C DETERMINE THE NEW VALUE OF THE VARIABLE, INDICATE IF SOLUTION IS
C INFEASIBLE SO THAT A NEW PATH CAN BE TAKEN, AND DEFINE THE NEW COLUMN
C TO BE ADDED TO THE INITIAL TABLEAU
430 CALL NZERO(N,NN,II,IN,ITSV,MV,E,KL,IA,BA,KM)
  IF(BA.NE.0.0) GO TO 430
C IF COLOR CHOSEN IS AT MAXIMUM VALUE, GO TO ANOTHER COLUMN
  IF(E.NE.0.0) CALL KLRMCH(N,NN,II,AD,MV,NZ,IN,ITSV,IA,E,MT,KL,KOLSV,
1KNT,KM,NC,IND,NU,N3,NX3)
C IF NO OTHER COLUMN CAN BE USED, RETURN TO NEXT FEASIBLE ITERATION
  IF(AD.NE.0.0) GO TO 411
C SAVE NECESSARY VALUES TO BE USED IN RECONSTRUCTING THIS ITERATION IF
441 DO 440 I=1,N

```

IBAK0114
IBAK0115

IBAK0116
IBAK0117
PREVIOUS

IBAK0118
IBAK0120
IBAK0121
IBAK0125
IBAK0130
IBAK0132
IBAK0133

IBAK0157
IBAK0158

IBAK0160
IBAK0165
IBAK0170
IBAK0180
IBAK0190
IBAK0210
IBAK0220
IBAK0224
IBAK0228
IBAK0230
RETURN

IBAK0235
IBAK0240

IBAK0250

IBAK0251
IBAK0252

IBAK0253
IBAK0254

IBAK0255
REQUIRED
IBAK0258

```

440 KOLSV(I,KNT)=IA(I)
    KOLSV(N+1,KNT)=II
    KOLSV(N+3,KNT)=MV(II)
    RETURN
450 D=NN
    WRITE(6,451)
451 FORMAT(10X,'ALL COMBINATIONS HAVE BEEN TRIED, NO SOLUTION EXISTS.')
```

IBAK0280
IBAK0290
IBAK0305
IBAK0310
IBAK0320
IBAK0330
IBAK0340
IBAK0350

```

    RETURN
    END
```

SUBROUTINE KVAR(NN,MZ,II,AD,MV,ISV,NC,IND,NU,KNT,NX3) KVAR0010
SUBPROGRAM TO FIND THE NEXT COLUMN IN LINE WITH THE LEAST NUMBER OF ZEROS.
C SYMBOLS ARE: MZ=NZ ; OTHERS AS PREVIOUSLY DEFINED

C DIMENSION MZ(NN),MV(NN),ISV(NN),NC(NN),NU(NX3) KVAR0020
C ASSIGN INITIAL VALUES
ISV(II)=0
IF(NU(KNT).NE.0) GO TO 501
NP=NN-1
C IF ALL COLUMNS WITH NONBORDERING ZEROS HAVE BEEN EXAMINED, EXAMINE TOTAL
C NUMBER OF ZEROS IN EACH COLUMN FOR NEXT COLUMN TO USE.
IF(IND.NE.0) GO TO 540
C CHECK THE REST OF THE COLUMNS IN THE MATRIX FOR THE NEXT COLUMN WITH
C NUMBER OF NONBORDERING ZEROS.
DO 520 K=II,NP
IF(NC(K+1).EQ.NC(II)) GO TO 530
520 CONTINUE
C IF NO OTHER NONBORDERING ZERO COLUMN CAN BE CHOSEN, ASSIGN POSITIVE VALUE
C TO INDICATOR
IND=NN
II=1
C DETERMINE WHICH COLUMN HAS THE LEAST TOTAL NUMBER OF ZEROS
DO 510 K=1,NP
510 IF(MZ(K+1).LT.MZ(II)) II=K+1
C CHECK THE REST OF THE COLUMNS IN THE MATRIX FOR THE NEXT COLUMN WITH
C TOTAL NUMBER OF ZEROS.
540 DO 500 K=II,NP
IF(MZ(K+1).EQ.MZ(II)) GO TO 530
500 CONTINUE
C IF NO OTHER COLUMN WITH SAME NUMBER OF ZEROS EXISTS, NOTE SAME AND RETURN
501 AD=NN
RETURN
C DEFINE NEW VARIABLE TO BE USED
530 II=K+1
RETURN
END

KVAR0030
KVAR0035
KVAR0040
KVAR0048
KVAR0050
KVAR0051
KVAR0052
KVAR0053
KVAR0054
KVAR0057
KVAR0058
KVAR0059
KVAR0060
KVAR0070
KVAR0080
KVAR0090
KVAR0100
KVAR0130


```

C SUBROUTINE INFEAS(N,NN,MV,IN,II,J,KL,IA,BA,KM,K2) INFS0010
C AS PREVIOUSLY DEFINED SYMBOLS ARE
C
C DIMENSION MV(NN),IN(N,NN),IA(N),KM(NN) INFS0020
C INCREASE COUNTER BY ONE; INDICATES WHEN ALL BORDERS HAVE BEEN CHECKED
  K2=K2+1 NZER0033
C EXAMINE ROW FROM STATE#1 TO CHOSEN STATE MINUS 1 TO SEE IF AN ADJACENT
  STATE IS DEFINED IN THIS AREA INFS0040
  IL=II-1
  DO 720 KK=1,IL INFS0050
    C IF THERE EXISTS AN ADJACENT AREA, CHECK THE TWO COLORS FOR A MATCH
    IF(IN(J,KK).NE.0) CALL ICHK(N,NN,II,IN,MV,KK,KL,IA,BA,KM,K2) INFS0060
    IF(BA.NE.0.0) RETURN INFS0070
  720 CONTINUE INFS0075
  C EXAMINE ROW FROM CHOSEN STATE PLUS 1 TO STATE#NN TO SEE IF AN ADJACENT
  STATE IS DEFINED IN THIS AREA INFS0090
  IU=II+1 INFS0100
  DO 725 KK=IU,NN INFS0110
    C IF THERE EXISTS AN ADJACENT AREA, CHECK THE TWO COLORS FOR A MATCH
    IF(IN(J,KK).NE.0) CALL ICHK(N,NN,II,IN,MV,KK,KL,IA,BA,KM,K2) INFS0120
    IF(BA.NE.0.0) RETURN INFS0130
  725 CONTINUE INFS0135
  RETURN INFS0150
  END INFS0160

C SUBROUTINE ICHK(N,NN,II,IN,MV,K,KL,IA,BA,KM,K2) ICHK0010
C THE ROUTINE RETURNS TO FIND ANOTHER COLOR OR VARIABLE IF THE COLORS ARE THE SAME
C WHEN THE COLOR HAS BEEN ESTABLISHED, THE NEW COLUMN TO ADD TO THE INITIAL
C TABLEAU IS DEFINED. SYMBOLS ARE AS PREVIOUSLY DEFINED.
C
C DIMENSION MV(NN),IA(N),IN(N,NN),KM(NN) ICHK0020
C
C IF ADJACENT AREAS ARE THE SAME COLOR, RETURN TO STARTING POINT AND REDEFINE
  COLOR OR VARIABLE
  IF(MV(II).EQ.MV(K)) GO TO 830 ICHK0040
C WHEN ALL OF AN AREA'S BORDERS HAVE BEEN CHECKED, DEFINE THE NEW COLUMN TO
  BE ADDED
  IF(KM(II).EQ.K2) GO TO 813 ICHK0056
C IF ALL BORDER AREAS HAVE NOT BEEN CHECKED, RETURN TO SUBROUTINE INFEAS
  AND CHECK NEXT BORDER AREA
  RETURN ICHK0057

```

```

C      813 DO 820 J=1,N
      820 IA(J)=0
      DO 825 J=1,N
      DO 825 JJ=1,NN
      IA(J)=IA(J)+IN(J,JJ)*MV(JJ)
      825 CONTINUE
      RETURN
      830 BA=N
      RETURN
      END

      SUBROUTINE KLRMCH(N,NN,II,AD,MV,NZ,IN,ITSV,IA,E,MT,KL,KSV,KT,KM,
      INC,IND,NU,N3,NX3)
      SUBROUTINE TO DEFINE ANOTHER COLUMN TO USE, IF THE SOLUTION IS INFEASIBLE
      AND AT MAXIMUM COLOR VALUE. SYMBOLS ARE: AS DEFINED ELSEWHERE

      DIMENSION MV(NN),NZ(NN),IN(N,NN),ITSV(NN),IA(N),MT(N,NN),KM(NN)
      DIMENSION NC(NN),KSV(N3,NX3),NU(NX3)

      C      ZERO DUMMIES USED
      700 E=0.0
      BA=0.0
      KE=KT-1
      C      DEFINE COLOR OF VARIABLE TO BE A PREVIOUS COLOR USED.
      MV(II)=0
      DO 740 I=1,KE
      IF(KSV(N+1,I)).EQ.II) MV(II)=KSV(N+3,I)
      740 USE THE NEXT COLUMN WITH SAME NUMBER OF ZEROS. IF NO SUCH COLUMN EXISTS,
      RETURN TO NEXT APPLICABLE ITERATION
      CALL KVAR(NN,NZ,II,AD,MV,ITSV,NC,IND,NU,KT,NX3)
      C      SAVE THE VALUE OF THE INDICATOR
      KSV(N+2,KT)=IND
      IF(AD.NE.0.0) RETURN OF THE VARIABLE, INDICATE IF SOLUTION IS
      C      DETERMINE THE NEW VALUE OF THE VARIABLE, INDICATE IF SOLUTION IS
      C      INFEASIBLE SO THAT A NEW PATH CAN BE TAKEN, AND DEFINE THE NEW COLUMN
      C      TO BE ADDED TO THE INITIAL TABLEAU
      710 CALL NZERO(N,NN,II,IN,ITSV,MV,E,KL,IA,BA,KM)
      IF(BA.NE.0.0) GO TO 710
      IF(E.NE.0.0) GO TO 700
      C      REZERO PRESENT INDICATOR
      IND=0
      RETURN
      END

      SUBROUTINE KONFLT(NN,J1,K1,NZ,NZCON,IN,N)

```

ICHK00068
 ICHK00070
 ICHK00080
 ICHK00090
 ICHK00100
 ICHK00110
 ICHK00120
 ICHK00130
 ICHK00140

KLMH00010
 KLMH00011
 KLMH00012
 KLMH00013
 KLMH00014

KLMH00020
 KLMH00021

KLMH00030
 KLMH00031
 KLMH00040

KLMH00050
 KLMH00051
 KLMH00052

KLMH00060
 KLMH00065
 KLMH00070

KLMH00080
 KLMH00085
 KLMH00090

KLMH00100
 KLMH00120

KNFT00010

```

C SUBPROGRAM TO COUNT THE NUMBER OF ZEROS IN EACH COLUMN AND TO COUNT THE
C NUMBER OF ZEROS ASSOCIATED WITH NONBORDERING AREAS. SYMBOLS ARE AS
C PREVIOUSLY DEFINED
C
  DIMENSION NZ(NN),NZCON(NN),IN(N,NN)
  NZ(J1)=NZ(J1)+1
  IF(IN(K1,J1).EQ.0) NZCON(J1)=NZCON(J1)+1
  RETURN
  END
  KNFT0020
  KNFT0030
  KNFT0040
  KNFT0050

C SUBROUTINE NOMCH(N,NN,IIS,II,KK,MV,IA,IN,NU,KNT,KL,ITSV,NX3)
C IN THE EVENT THAT THE COLOR CHOSEN WHEN ENUMERATING ALL POSSIBLE COLORS
C FOR THE CHOSEN AREA DOES NOT MATCH ANY OF THE BORDERING AREAS, THIS
C SUBPROGRAM DEFINES THE VARIABLE, ITS COLOR, AND THE NEW COLUMN TO ADD
C
  DIMENSION MV(NN),IA(N),IN(N,NN),ITSV(NN),NU(NX3)
  C DEFINE CHCSEN VARIABLE
  II=IIS
  C DEFINE VARIABLE'S COLOR
  MV(II)=KK
  C IF COLOR IS AT THE MAXIMUM VALUE, MAKE A NOTE OF THE VARIABLE
  IF(MV(II).EQ.KL) ITSV(II)=N
  C DEFINE THE COLUMN TO ADD
  DO 1200 J=1,N
  1200 IA(J)=0
  DO 1251 J=1,N
  1251 JJ=1,NN
  DO 1250 IA(J)=IA(J)+IN(J,JJ)*MV(JJ)
  1250 CONTINUE
  KNT=KNT+1
  NU(KNT)=N
  RETURN
  END
  NMCH0020
  NMCH0030
  NMCH0040
  NMCH0050
  NMCH0060
  NMCH0070
  NMCH0080
  NMCH0090
  NMCH0100
  NMCH0110
  NMCH0115
  NMCH0118
  NMCH0120

C SUBROUTINE IORDER(N,KN1,ITER)
C SUBPROGRAM TO ORDER THE ITERATIONS CHOSEN FOR BACKTRACKING
C SYMBOLS ARE AS DEFINED IN SUBROUTINE IBAK
C
  DIMENSION ITER(N)
  KN1M1=KN1-1
  DO 1010 J=1,KN1M1
  IG=ITER(J)
  J1=J+1
  DO 1011 I=J1,KN1
  IF(IG.GE.ITER(I)) GO TO 1011
  IG=ITER(I)
  IORD0010
  IORD0020
  IORD0030
  IORD0035
  IORD0036
  IORD0037
  IORD0038
  IORD0039
  IORD0040

```


IORD00041
 IORD00042
 IORD00043
 IORD00070
 IORD00080

ITER(J)=ITER(J)
 ITER(J)=IG
 1011 CONTINUE
 1010 RETURN
 END

C SUBROUTINE LKINF(N,NN,IIS,KNI,IN,MV,MVMCH,JJ)
 C SUBROUTINE TO DETERMINE WHICH STATES BORDER ON CHOSEN STATE. SYMBOLS
 C ARE AS PREVIOUSLY DEFINED LKIF0010

C DIMENSION IN(N,NN),MV(NN),MVMCH(NN)
 C IF(IIS.EQ.1) GO TO 910
 C EXAMINE ROW FROM STATE #1 TO CHOSEN STATE MINUS 1 TO SEE IF AN ADJACENT
 C STATE IS DEFINED IN THIS AREA.

C ISL=IIS-1
 C DO 920 KK=1,ISL
 C IF THERE EXISTS AN ADJACENT AREA, CHECK THE TWO COLORS FOR A MATCH
 C IF(IN(JJ,KK).NE.O) CALL LKICHK(NN,IIS,KK,KNI,MV,MVMCH)

920 CONTINUE
 C IF(IIS.EQ.NN) RETURN
 C EXAMINE ROW FROM CHOSEN STATE PLUS 1 TO STATE #NN TO SEE IF AN ADJACENT
 C STATE IS DEFINED IN THIS AREA.

C 910 ISU=IIS+1
 C DO 925 KK=ISU,NN
 C IF THERE EXISTS AN ADJACENT AREA, CHECK THE TWO COLORS FOR A MATCH
 C IF(IN(JJ,KK).NE.O) CALL LKICHK(NN,IIS,KK,KNI,MV,MVMCH)

925 CONTINUE
 RETURN
 END

C SUBROUTINE LKICHK(NN,IIS,K,KNI,MV,MVMCH)
 C SUBROUTINE COMPARES THE COLORS OF ADJACENT AREAS. IF THE COLORS ARE
 C SAME, THE ROUTINE RECORDS WHICH AREA HAS SAME COLOR AS CHOSEN
 C AREA. SYMBOLS ARE AS DEFINED IN IBAK LKIK0010

C DIMENSION MV(NN),MVMCH(NN)
 C COMPARE THE COLORS OF THE ADJACENT AREAS
 C IF(MV(IIS).EQ.MV(K)) GO TO 890

C RETURN
 C RECORD THE VARIABLE WITH THE MATCHING COLOR
 C 890 KNI=KNI+1
 C MVMCH(KNI)=K
 C RETURN
 C END

SUBROUTINE KNTVAR(N,NN,IN,I,IITMP,L2) KNV00010

LKIK0020
 LKIK0030
 LKIK0040
 LKIK0044
 LKIK0045
 LKIK0050

```

C SUBPROGRAM TO DETERMINE WHICH TWO VARIABLES ARE ADJACENT AND STILL HAVE
C COLOR ZERO.  SYMBOLS ARE AS DEFINED IN IBAK
C
      DIMENSION IN(N,NN),IITMP(NN)
      DO 1300 J=1,NN
      IF (IN(I,J).NE.0) L2=L2+1
      1300 IF (IN(I,J).NE.0) IITMP(L2)=J
      RETURN
      END
      KNVR00020
      KNVR00030
      KNVR00040
      KNVR00050
      KNVR00060

```

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KEY WORDS

LINK A

LINK B

LINK C

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Four color problem





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An iterative process to solve the graph-



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